

IBM System z

## IBM z6 – The Next-Generation Mainframe Microprocessor

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## IBM z6 Highlights

 New high-frequency z/Architecture microprocessor core

>4 GHz operation in system

- 4 cores per die
  - Each with 3MB private 2<sup>nd</sup>-level cache
- Accelerator engines
  - Data compression
  - Cryptographic functions
  - Decimal floating point
- Integrated SMP communications
  - Switch connects cores to SMP Hub chip
    - Shared cache and SMP fabric
  - Memory bus controller
  - I/O bus controller
  - El3 technology up to 3 GHz bus speeds
- System interfaces
  - 2 x 48 GB/s SMP Hub
  - 4 x 13 GB/s Memory
  - 2 x 17 GB/s I/O

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- 991M Transistors
- 138 Mb SRAM
- 6 km wire
- 21.7 X 20.0 mm die
- 1188 signal / 8765 total chip I/Os

#### **Relationship to Power6**

- Siblings, not identical twins
- Share lots of DNA
  - -IBM 65nm SOI technology
  - -Design building blocks:
    - latches, SRAMs, regfiles, dataflow elements
  - -Large portions of FXU, BFU, DFU, MC, GX
  - -EI3 interface technology
  - -Core pipeline design style
    - •High-frequency, low-latency, mostly-in-order
  - -Many designers
- Different personalities
  - -Very different ISAs => very different cores
  - -Cache hierarchy and coherency model
  - -SMP topology and protocol
  - -Chip organization
  - -IBM z6 optimized for Enterprise Data Serving Hub





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#### IBM z6 Instruction Set Architecture

- Continues line of upward-compatible mainframe processors
  - Application compatibility since 1964
  - Supports all z/Architecture-compliant OSes



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### IBM z6 Architecture

Continues line of upward-compatible mainframe processors

- Rich CISC ISA
  - 894 instructions (668 implemented entirely in hardware)
  - 24, 31, and 64-bit addressing modes
  - Multiple address spaces robust inter-process security
  - Multiple arithmetic formats
  - Industry-leading virtualization support
    - High-performance logical partitioning via PR/SM
    - Fine-grained virtualization via zVM scales to 1000s of images
  - Precise, model-independent definition of hardware / software interface
- Architectural extensions for IBM z6
  - 50+ instructions added to improve compiled code efficiency
  - Enablement for software/hardware cache optimization
  - Support for 1MB page frames
  - Full hardware support for Decimal Floating Point



#### IBM z6 Continues the CMOS Mainframe Heritage



G4 – 1<sup>st</sup> full-custom CMOS S/390
 G5 – IEEE-standard BFP; branch target prediction
 G6 – Cu BEOL

z900 – Full 64-bit z/Architecture
 z990 – Superscalar CISC pipeline
 z9 EC – System level scaling

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#### **IBM z6 Microprocessor Core**



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#### **High-Frequency Design**

- Frequency-optimized methodology
  - -Logic / circuit co-design
  - Reduced latch delay overhead

#### -Arithmetic dataflow stacks shared with Power6

- Minimize critical pipeline latencies
- Isolate hardware complexity of CISC ISA
- Optimization of cache lookup path
- Aggressive branch prediction





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### **High-Frequency Design**

- Frequency-optimized methodology
- Minimize critical pipeline latencies
  - Overall pipeline depth increased significantly from predecessor
    - •15FO4 vs. 28FO4 cycle
  - Minimal increase in performance-critical core loop latency
  - ISA complexity pushed to front end of pipeline
    - outside of key dependency loops
  - -Non-stalling pipeline:
    - •once issued, instruction either finishes or is recycled
- Isolate hardware complexity of CISC ISA
- Optimization of cache lookup path
- Aggressive branch prediction



#### **Core Pipeline**



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#### **Core Pipeline**



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|               | Cycles |        | FO4    |        |
|---------------|--------|--------|--------|--------|
| Pipeline Loop | IBM z9 | IBM z6 | IBM z9 | IBM z6 |
| Load-FXU      | 0      | 0      | 0      | 0      |
| FXU-FXU       | 1      | 1      | 28     | 15     |
| Load-Load     | 3      | 4      | 84     | 60     |
| FXU-Load      | 4      | 5      | 112    | 75     |
| Branch wrong  | 6+     | 13+    | 168+   | 195+   |



Decode

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### **High-Frequency Design**

- Frequency-optimized methodology
- Minimize critical pipeline latencies
- Isolate hardware complexity of CISC ISA
  - Multi-pass handling of special cases
  - Leverage millicode for complex operations
- Intense optimization of Address\_Add / Cache lookup / FXU path
  - Cache set prediction; latch boundary adjustment
  - Performance / array / logic co-optimization
    - →128KB D-Cache and 4-cycle loop
- Aggressive branch prediction





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### **High-Frequency Design**

- Frequency-optimized methodology
- Minimize critical pipeline latencies
- Isolate hardware complexity of CISC ISA
- Optimization of cache lookup path
- Aggressive branch prediction
  - Minimize impact of long front end of pipe



## **Branch Prediction**

- Multiple history-based prediction mechanisms
  - 2 level Branch Target Buffer
  - Filtered Pattern History Table
  - Tagged multi-target prediction
  - Level 2 BTB data compression





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## **Branch Prediction**

- Multiple history-based prediction mechanisms
  - 2 level Branch Target Buffer
  - Filtered Pattern History Table
  - Tagged multi-target prediction
  - Level 2 BTB data compression
- BTB searches run ahead of I-fetch
  - Enables I-cache prefetching of targets
  - Prevents fall-through false misses in I-cache
- Tight interaction with I-Fetch and IDU
  - Basic block merging
  - Relative branch snooping
- z/Architecture Optimizations
  - Millicode entry / exit prediction
  - EXecute instruction optimization
  - Efficient handling of "indirect" branches



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#### **Decimal Floating Point Accelerator**

- Meets requirements of business and human-centric applications
  - Performance, Precision, Function
  - Avoids rounding and other problems with binary/decimal conversions
  - Improved numeric functionality over legacy BCD operations
  - Much of commercial computing is dominated by decimal data and decimal operations





144 bits or 36 digits wide

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### **Decimal Floating Point Accelerator**

Meets requirements of business human-centric applications
 Performance, precision, Function

- IBM z6 DFU co-developed with Power6
  - -Common architecture operations and semantics
  - -Common dataflow elements
  - -Mainframe legacy BCD operations mapped onto DFU in IBM z6



144 bits or 36 digits wide

- Growing industry support for DFP standardization
  - -Java BigDecimal,C#, XML, C/C++, GCC, COBOL all supporting IEEE 754R
  - -Endorsed by key software vendors including Microsoft and SAP
  - -Open standard definition led by IBM



#### **Compression and Cryptography Accelerator**

- Data compression engine
- Cryptography engine
- Accelerator unit shared by 2 cores





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#### **Compression and Cryptography Accelerator**

- Data compression engine
  - Static dictionary compression and expansion
  - Dictionary size up to 64KB (8K entries)
    - Local 16KB caches for dictionary data
  - Up to 8.8 GB/sec expansion
  - Up to 240 MB/sec compression
- Cryptography engine
  - DES (DEA, TDEA2, TDEA3)
  - SHA-1 (160 bit)
  - SHA-2 (256, 384, 512 bit)
  - AES (128, 192, 256 bit)
  - 290-960 MB/sec bulk encryption rate
- Accelerator unit shared by 2 cores
  - Independent compression engines
  - Shared cryptography engines
  - Co-operates with core millicode
  - Direct path into core store buffers





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#### Industry-Leading Error Detection and Recovery

- Fine-grained redundancy and checking throughout design
  - –ECC on 2<sup>nd</sup> and 3<sup>rd</sup>-level caches, store buffers, R-Unit state array
  - -Parity on all other arrays, register files
  - -Parity or residue checking on data/address/execution flow
  - -Extensive functional, parity, and state checking on control logic
  - -Over 20,000 error checkers in chip
- Full architected state of processor buffered in R-Unit with ECC
  - -Allows precise core retry for almost all hardware errors
  - -Dynamic, transparent core sparing in the event of hard error in core
- Machine check architecture allows precise software recovery

   Minimizes system impact in rare case of unrecoverable failure



#### Industry-Leading Error Detection and Recovery

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Machine check architecture allows precise software recovery

 Minimizes system impact in rare case of unrecoverable failure
 Target is Bullet-proof computing:
 Never lose data

✓Never go down

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## SMP Hub Chip

- Connects multiple IBM z6 processor chips
   –48 GB/Sec bandwidth per processor
- Shared 3<sup>rd</sup>-Level cache
  - -24MB SRAM Cache
  - -Extended directory
    - Partial-inclusive discipline
  - -Hub chips can be paired
    - 48MB shared cache
- Low-latency SMP coherence fabric –Robust SMP scaling
  - -Strongly-ordered architecture
- Multiple hub chips/pairs allow further SMP scaling



- 1.6 Billion Transistors
- 242 Mb SRAM
- 3 km wire
- 20.8 X 21.4 mm die
- 2419 signal / 7984 total chip I/Os

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#### **Energy Efficiency**

- Mainframe focus on system and data center efficiency
  - Consolidation of many servers onto one system
  - **Consistent performance at sustained high utilization** 
    - Resilience in face of changing workloads
  - Leverages virtualization capabilities of PR/SM, zVM, zOS
- System z designs are optimized for scale-up data serving
  - SMP Hub design enables robust scaling across wide spectrum of workloads
  - Centralized SMP fabric minimizes fabric logic per core
    - Extended on IBM z6 via 4-core processor chip
  - MRU cooling allows dense package and reduces leakage power
  - Extensive hardware support for multi-level virtualization
- Chip-level power optimization applied to IBM z6 design
  - Local clock gating to limit maximum dynamic power
  - Millicode sleep mode for wait/spare/stop states



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#### Conclusion

- IBM z6 processor takes mainframe computing to the next level
  - New high-frequency 4-core microprocessor
  - Advanced pipeline design
  - Enhanced performance on CPU-intensive workloads
- IBM z6 is specifically designed and optimized for mainframe systems
  - Full z/Architecture compatibility
  - New features enhance enterprise data serving performance
  - Industry-leading virtualization capabilities
  - Energy efficiency at system and data center levels
  - Extends heritage of industry-leading RAS

# Thank You!

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